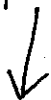


# INTRODUCTION TO PROOF

①

- Theorem :** A statement that can be shown to be true
- Axioms :** Statements that we assume to be true
- Lemma :** A less important theorem that is helpful in the proof of other results
- Corollary :** It is a theorem that can be established directly from a proven theorem.

**Conjecture :** Conjecture is a statement that is proposed as true statement, usually on the basis of some partial evidence.



A heuristic argument.

☐ When a conjecture is proved, it becomes a theorem.

☐ **Rule of thumb:**

To prove a universal statement, it must be shown that it works for all cases

To disprove a universal statement, one counter example is enough.



# MATHEMATICAL PROOFS ...

Direct Proof :

Direct proof is a mathematical argument that uses rules of inference to draw conclusion out of the premises.

Let's consider Disjunctive syllogism

$$\begin{array}{r}
 \downarrow \\
 P \vee Q \quad \text{Premise 1} \\
 \neg P \quad \text{Premise 2} \\
 \hline
 \end{array}$$

$\therefore Q$

We can use direct proof method, that is, a chain of inferences,

$P \vee Q$	Premise
$Q \vee P$	Commutivity of $\vee$
$\neg(\neg Q) \vee P$	Double negation law
$\neg Q \rightarrow P$	$A \rightarrow B \equiv \neg A \vee B$
$\neg P$	Premise 2
$\neg\neg Q$	Modus Tollens
<hr/>	
$Q$	Conclusion

Generally, when we want to prove a conditional statement  $P \rightarrow Q$ , we assume "P" as true and follow implications to show that Q is true as well.

↓ Direct proof

We need to find the propositions that obtain Q as the conclusion.

□ Prove: Given  $m$  is even and  $n$  is odd, their sum  $(m+n)$  is always odd.

By definition of odd and even:

If there's an integer  $j$ , then

odd " $n$ " = $2j+1$	Their sum,	$m+n = (2j+1) + 2k$
even " $m$ " = $2k$		

↳ an integer

So,  $(j+k)$  is an integer. Thus,  $m+n$  is odd by definition.

This is direct proof.

□ Comments on direct proofs:

- In direct proofs
- ✦ We start with hypothesis
  - ✦ Continue with a sequence of deductions
- ↓ end with  
Conclusion

However, this may not be true always. We may reach to dead ends if direct methods are followed.

\*  $n$  is an even integer,  $n^2$  is even

Say,  $n = 2k$ . Then,

$n^2 = 4k^2$

⇒  $n^2 = 2 \cdot (2k^2)$

$= 2 \cdot \text{integer}$

$\equiv \text{even}$

Proved

We need indirect methods that do not start with hypothesis.

Contrapositive:

" If it is <sup>P</sup>hartal today,  
then I do not go to class"  
<sup>Q</sup>

↓ Contrapositive

" If I go to class, then  
it is not <sup>¬Q</sup>hartal today"  
<sup>¬P</sup>

Converse:

" If I do not go to class, then it is  
hartal today "

Not true; fallacy of the converse.

□ Given, 'n' is an integer and  $3n+2$  odd, then 'n' is odd.

Direct Proof:

$$3n+2 = 2k+1$$

$$\Rightarrow 3n = 2k-1$$

what's next?

No direct way to proceed further

Indirect Method: Proof by contraposition

So, according to contraposition theorem

we assume n is even

↓ and we show

$3n+2$  is even

□ So,  $n = 2k$ , for some integer 'k'

$$3n+2 = 3 \cdot 2k + 2 = 2(3k+1)$$

this is another integer

≡ even

□ Proof by contraposition continues ...

Because the negation of the conclusion is false and it implies that the hypothesis false, therefore the original conditional statement is true

□ Another example

Assume  $x \in \mathbb{Z}$ . Prove that  $x^2 - 6x + 5$  is even, then  $x$  is odd.

Let's consider

$$x = 2a \text{ for any integer "a"}$$

↳ we start with  $\neg q$ ; contrapositive  
↳ even.

So, we obtain,

$$\begin{aligned}
x^2 - 6x + 5 &= (2a)^2 - 6 \cdot 2a + 5 \\
&= 4a^2 - 12a + 5 \\
&= 4a^2 - 12a + 4 + 1 \\
&= 2(2a^2 - 6a + 2) + 1 \\
&= 2k + 1 \\
&\equiv \text{Odd}
\end{aligned}$$

So, we started with  $\neg q$  and we show that

$$x^2 - 6x + 5 \text{ is odd; that is } \neg p$$

That's, we prove that  $x^2 - 6x + 5$  is odd.

□ Proof by contradiction

In this proof method, we assume that the statement made is not true.

↓ then  
we derive a contradiction

□ Example Prove that  $\sqrt{2}$  is irrational.  
we assume rational

Let's assume  $\sqrt{2}$  is rational

↓ implies,  $\sqrt{2}$  can be written as  $m/n$ , where  $m, n$  are integers

↓  
 $m^r/n^r = 2 \Rightarrow m^r = 2n^r \Rightarrow m^r$  is even  
 $\Rightarrow m$  is even  
 $\Rightarrow m = 2k$ .

Now,

$(2k)^r = 2n^r$   
 $\Rightarrow 2n^r = 4k^r$   
 $\Rightarrow n^r = 2k^r \Rightarrow n$  is also even.

Thus,  $\forall m$  and  $n$  are both even and they have a common factor 2.

$\forall$  This contradicts the assumption that  $m/n$  was in lowest terms

↳ Not in lowest term

So, by contradiction, it can be concluded that  $\sqrt{2}$  is irrational.

# Proof by Cases

$p \rightarrow r$	Premise 1
$q \rightarrow r$	Premise 2
$p \vee q$	Premise 3
$\therefore r$	conclusion

Example: Given  $x$  is an integer  
 $x^2 + x$  is even

Set-up for proof-by-cases:

Let's assume  $p: x$  is even |  $r: x^2 + x$  is even  
 $q: x$  is odd

Verify Premise 1: if  $x$  is even,  
 $x = 2n$  |  $x^2 + x = (2n)^2 + 2n$   
 $= 4n^2 + 2n$   
 $= (2n)^2 + 2n$

Verify Premise 2: if  $x$  is odd  
 $x = 2n+1$  |  $x^2 + x = (2n+1)^2 + 2n+1$  ↳ which is even  
 $= (4n^2 + 4n + 1) + 2n + 1$   
 $= 4n^2 + 6n + 2$

Verify Premise 3: An arbitrary integer is ↳ even  
either even or odd.

So, the conclusion is proved.

~~□~~ Mistakes in Proof

⇒  $a = b$  Given

⇒  $a^2 = ab$  Multiplied by  $a$

⇒  $a^2 - b^2 = ab - b^2$  subtract  $b^2$

⇒  $(a-b)(a+b) = b(a-b)$

⇒  $(a-b)(a+b) = b(a-b)$  Mistake is here; divide by  $(a-b)$

⇒  $a + b = b$

⇒  $2b = b$  Replace  $a$  by  $b$  as  $a = b$

⇒  $2 = 1$  | Where's the mistake?  
0

□ Proof by Cases: another example

Let  $n$  be an integer. show that if  $n$  is not divisible by 3, then  $n^2 = 3k + 1$  for some integer  $k$

Assume  
Case 1:  $n = 3m + 1$ . So,  $n^2 = (3m + 1)^2 = 9m^2 + 6m + 1 = 3(3m^2 + 2m) + 1$   
Not divisible by 3 =  $3(\underbrace{3m^2 + 2m}_{\text{integer}}) + 1 = 3k + 1$

Case 2:  $n = 3m + 2$ . So,  $n^2 = (3m + 2)^2 = 9m^2 + 12m + 4$   
 $= 9m^2 + 12m + 3 + 1$   
 $= 3(\underbrace{3m^2 + 4m + 1}_{\text{integer } k}) + 1$   
 $= 3k + 1$ , which is not divisible by 3

So, Case I and Case II reflect all possible possibilities. Thus, proved.



## □ Proofs of Equivalence

⑦

To prove a theorem that is a biconditional statement, that is, statement of the form  $P \leftrightarrow Q$

we show that

$P \rightarrow Q$  and  $Q \rightarrow P$  are true

This approach is based on:

$$(P \leftrightarrow Q) \leftrightarrow [(P \rightarrow Q) \wedge (Q \rightarrow P)]$$

is the tautology.

For example: Given a theorem

If  $n$  is a positive integer, then  $n$  is odd if and only if  $n^2$  is odd

To prove this, we must show that

$$P \rightarrow Q, \quad Q \rightarrow P$$

where,

$P$ : " $n$  is odd"

$Q$ : " $n^2$  is odd"

## ☐ NORMAL FORM

A product of the variables and their negations in a formula is called an elementary product.

$\neg p \wedge q, q \wedge r$  are example of elementary products.

A sum of variables and their negations is called an elementary sum.

$\neg p \vee q, q \vee p \vee s$  are examples of elementary sum

☐ Elementary sum is the disjunction of literals.

Elementary product is the conjunction of literals.

Observation:

Necessary condition for an elementary product to be identically false is to have at least one pair of literals where one (p) is the negation ( $\neg p$ ) to generate the others.

$$\frac{p \wedge \neg p \wedge \dots \wedge q \dots}{F} \equiv F$$

For elementary sum,

it becomes a tautology if one pair exists where one is the negation of other.

$$\frac{P \vee \neg P}{T} \vee \dots \dots \dots \equiv T$$

### Disjunctive Normal Form:

It is the formula which is similar to, the original formula, But it consists of a sum of elementary product.

To translate any formula to disjunctive normal form — Replace  $\rightarrow$  and  $\leftrightarrow$  using  $\wedge$ ,  $\vee$ , and  $\neg$

Example:  $(P \rightarrow Q) \wedge \neg Q$  obtain DNF  
 $\equiv (\neg P \vee Q) \wedge \neg Q$   
 $\equiv (\neg P \wedge \neg Q) \vee (Q \wedge \neg Q)$

### Conjunctive Normal Form:

A formula which is equivalent to a given formula and consists of a product of elementary sums is called conjunctive normal form. CNF

Example:  $(P \rightarrow Q) \wedge \neg Q$   
 $\equiv (\neg P \vee Q) \wedge \neg Q \equiv \text{CNF}$

□ CNF continues...

(10)

∞ We can bring any formula to normal form

∞ conjunctive normal form is unique.

□ Example of CNF

$$\square (p \rightarrow q) \leftrightarrow (p \rightarrow r)$$

$$\equiv [(p \rightarrow q) \rightarrow (p \rightarrow r)] \wedge [(p \rightarrow r) \rightarrow (p \rightarrow q)]$$

$$\equiv [\neg(p \rightarrow q) \vee (p \rightarrow r)] \wedge [\neg(p \rightarrow r) \vee (p \rightarrow q)]$$

$$\equiv [\neg(\neg p \vee q) \vee (\neg p \vee r)] \wedge [\neg(\neg p \vee r) \vee (\neg p \vee q)]$$

$$\equiv [(p \wedge \neg q) \vee (\neg p \vee r)] \wedge [(p \wedge r) \vee (\neg p \vee q)]$$

$$\equiv [(p \wedge \neg q) \vee \neg p \vee r] \wedge [(p \wedge r) \vee (\neg p) \vee q]$$

$$\equiv [(p \wedge \neg q) \wedge (\neg q \vee \neg p) \vee r] \wedge [(p \wedge \neg p) \wedge (r \vee \neg p) \vee q]$$

$$\equiv (\neg q \vee \neg p \vee r) \wedge (\neg r \vee \neg p \vee q)$$