

History of Programming Languages.

Von Neumann provided us two important concepts :

- a) Shared-program technique
- b) Conditional Control transfer

Shared-program technique :

✘ Actual computer hardware should be simple and not need to be rewired everytime to perform the specified task

Instead,

✘ Complex program/instructions should be used to control the simple hardware.

↳ This allows faster reprogramming.

Conditional Control Transfer :

✘ This provides the notion of subroutine, or blocks of code that could be jumped to in any order.

instead of a single set of chronologically ordered steps

✘ Also suggests that code should be able to branch based on logical statements..

↳ if, else
↳ For loop

For instance,

If (expression)
THEN.

Loop such as with FOR

In general, conditional control transfer provides the idea of libraries

In c

scanf }
printf } stdio.h
 } headerfile

sqrt }
pow } Math.h

↘ block code that can be reused again and again.

↘ Generally, pre-compiled, where actual functionality is implemented.

↘ But, actual implementation of these functions are done in .lib file
Say, MATH.lib

So, library —

- a set of code that can be reused over and over
- Pre-compiled
- Available in standard form to be used in other code.

Programming Languages ...

Short code :

- Appeared at around 1949
 - First computer language to instruct electronic devices.
- { Programmer needed to change the statements into 1's and 0's
- First step towards the complex language of today.

In 1951, American scientist —

Grace Hopper
wrote the first compiler A (A-0)
First version

Assembly Language :

- Mnemonic symbols are used for instructions codes and memory locations
- A program called an assembler translates symbolic assembly language code to binary machine code.

Mnemonics & symbols :

Generally, mnemonics are something similar to abbreviation that helps to remember something.

For instance,

MOV stands for moving data between registers & memory

LD



Load data to given location.

ADD Add the numbers

• `.ORIG x3000` ; Address of the first instruction

LD R1, FIRST ^{memory location}

LD R2, SECOND

ADD R3, R2, R1 ; Add the number in R2 and R1, and place the sum in R3

ST R3, SUM ; Copy the number in R3 to memory location SUM.

Only one `.ORIG` per program
Where to start in placing things in memory.

... continues

HALT ; Halt the program

FIRST .FILL #5

SECOND .FILL #6

SUM .BLKW #

↳ Declare a group of characters in memory.

.END Tells where the program source ends.

FORTRAN :

↳ FORMula TRANslating System, generally popular among the mathematicians.

Shortcoming of Assembly languages :

↳ Lacks in abstraction capability of mathematical notation

int first = 5

int second = 6

int sum = first + second

sample C program

to add two numbers

↳ a notation or way of expressing ideas that makes them concise, simple, and easy for human mind to grasp.

- ✓ The other shortcomings of assembly language is —
 - it is often hardware specific; particular computer hardware architecture has its own machine language set.

Hence, customized dialect of assembly language is necessary.

☐ FORTRAN (FORMula TRANslating System)

- ✓ Designed and Developed by John Backus from IBM.
- ✓ Intended for Mathematicians and scientists
- ✓ Early versions were close to Assembly language
 - ↳ Later versions have undergone numerous revisions.
- ✓ It is still in use, and the versions are being updated.
 - ↳ in physical systems
 - ↳ Astrophysics

FORTAN → I → ... IV → 66 → 77
→ 90

☐ LISP (List Processing Language)

- ✘ Second-oldest high-level programming, widely used in artificial intelligence (AI).
- ✘ Designed by John McCarthy in 1958
- ✘ Basic and only data type was list; other data types were added later.

✘ LISP programs are written a set of lists

→ easy to modify and hence, grow on its own.

→ Formatting is different than syntax
standard boolean logic.

OR (x, y) parenthesized prefixed
used in LISP

x OR y standard Boolean logic

Data structure:

Pure lisp has two kind of data structure — i) Atoms ii) lists.

Atoms :

- Either symbols
 - ↳ identifiers
 - ↳ Numerical symbols/literals

→ atom
(A B C D)

Simple list, in which elements are restricted to atoms.

Lists :

- specified by delimiting their elements parentheses.
- Nested list structures are also specified by parentheses.

¹ (A ² (B C) ³ D ⁴ (E (F G))) (expression of the nested list)

Here,

1. atom A
2. sublist (B, C)
3. atom D

4. Sublist (E (F G))
↓
second element is another sublist.

☐ Lists are stored as single-linked list structures

has nodes

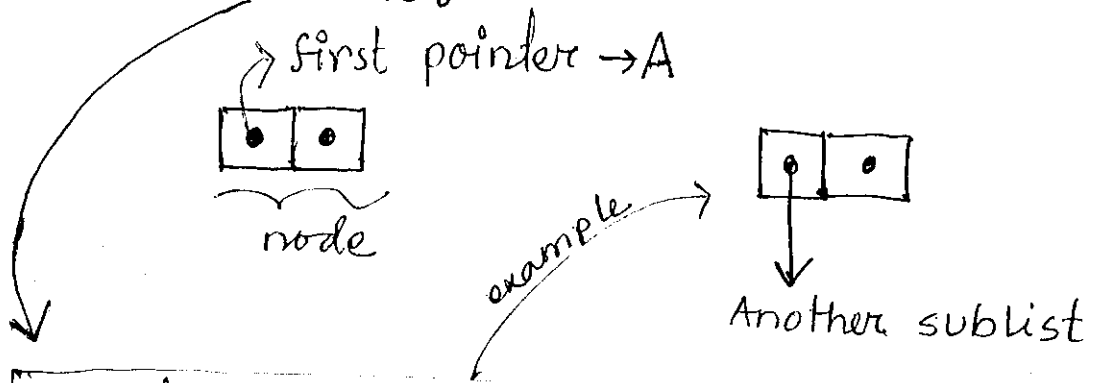
- ↳ each node has two pointers
- ↳ represents a list node.

✓ Node containing an atom has its first pointer pointing towards some representation of the atom.

- Symbol
- Numeric value

OR

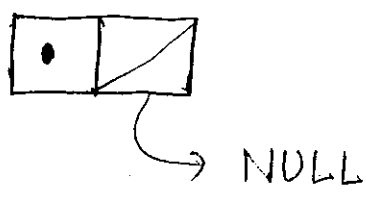
the first pointer may point toward a sublist



A node for sublist/element has its first pointer pointing towards the first node of the sublist

✓ In both cases, the second pointer of a node points to the next element in the list.

✓ Last element of a list doesn't have any successor — so, its link is NULL



Example: Internal representation of two Lisp lists.

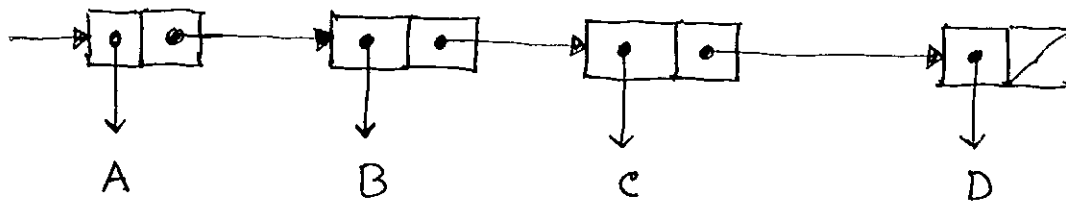


Fig. Representing list (A B C D)

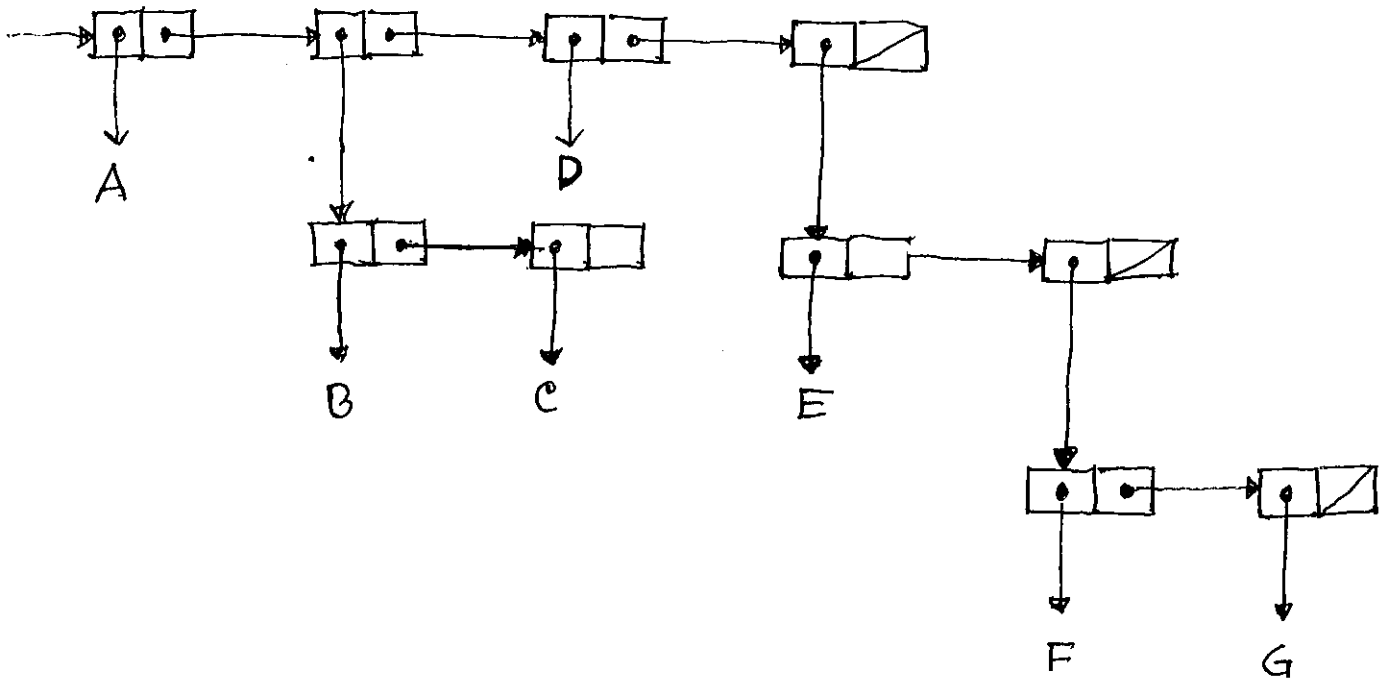


Fig. Representing list

(A (B C) D (E (F G)))

☐ More on LISP

- It is purely a functional programming

↙
↘
• Computations are performed by applying functions to arguments.

• Assignment statements and variables are not necessary in functional language program.

- Still the dominant language for AI

- Common LISP and Scheme are two main version.

- Symbolic Programming

→ Symbols: +, -, 1, 2, 3, 4 etc.

→ Symbolic Expression:
(+ 1 2), (+ (* 3 4) 2)

→ Generally, it manipulate symbolic expressions.

Descendants of LISP

Common LISP and Scheme are the two most widely used dialects.

SCHEME :

- Emerged from MIT in the mid 70s
- Small size
- Extensive use of static scoping

Also known as lexical scoping. Scoping provides the range of functionality of a variable

Because of this functionality range, the variable can be referenced from within the small block of code where the variable is defined.

scope is determined during compilation.

Example:

```
int x = 6  
main()
```

```
{  
  x = 3;
```

```
{  
  int x;
```

```
{  
  x = 2;
```

```
  printf("%d\n", x);  
}
```

```
  printf("%d\n", x);  
  return 0;  
}
```

Output: 2
 3

Static scoping

- Definition of a variable is set or resolved by
 - looking at its containing block
 - or, by looking at its function

↓ if fails
variable definition is resolved by searching the outer containing block
↓
and so on.

Example:

```
int a = 10, b = 20;
```

```
int main ()
```

```
{  
  int a = 5;
```

```
  {  
    int c;  
    c = b/a;  
    printf ("%d", c);  
  }
```

```
}
```

→ Containing block for variables a, b.

☐ ALGOL 58, came in 1958

✧ It was, in some sense, a descendant of Fortran

✧ Generalized many Fortran features and added several new constructs and concepts.

✧ Concept of data types was formalized

✧ Added the concept of compound statement. (begin ... end)

→ Other subsequent programming languages inherited this.

✧ Identifiers could be of any length,

Fortran I's restricted this to 6 or fewer characters

✧ Any dimensional array was allowed.

→ Fortran I's limitation was ≤ 3 dimensions.

✧ Nested selection of statements were possible

→ IS had an else-if clause was not in Fortran.

SCOPING

- ✓ Scoping is necessary to allow reuse of variables.
- ✓ An acceptable norm (and widely practiced as well) is to use shorter variable names. But in program that has millions of lines of code, shorter variable names may fall short.

↳ So, we must reuse variable names and scoping just allows it.

For instance,

```
int fnc 1 ( )  
{ int a = 5; }
```

scope of "a" within fn^c 1

```
int fnc 2 ( )  
{ int a = 10; }
```

Here, variable "a" is reused in fn^c 2. ↳
because of scoping
↳ is redefined here.
↳ scope of "a" is within fn^c

☐ Scoping Types

1. Static scoping
C, C++, Java etc.
2. Dynamic scoping
Perl (both static and dynamic)

Dynamic scoping

Let's consider a pseudocode :

```
{  
  printf ("%d", x)  
  ↪ for integer  
}
```

```
{  
  printf (x)  
}
```

We can use scoping concept to obtain the necessary information needed.

- Free variable
- Value not known
- data type not defined.

Assume,

```
Test ()  
{ print (x)  
}
```

```
Sample ()  
{  
  x = 20;  
  Test ()  
}
```

value of x is fixed at point of invocation.

Point of invocation.

so, dynamic scoping obtains $x=20$ and print accordingly

Whereas, for the above case, static scoping takes values from main()

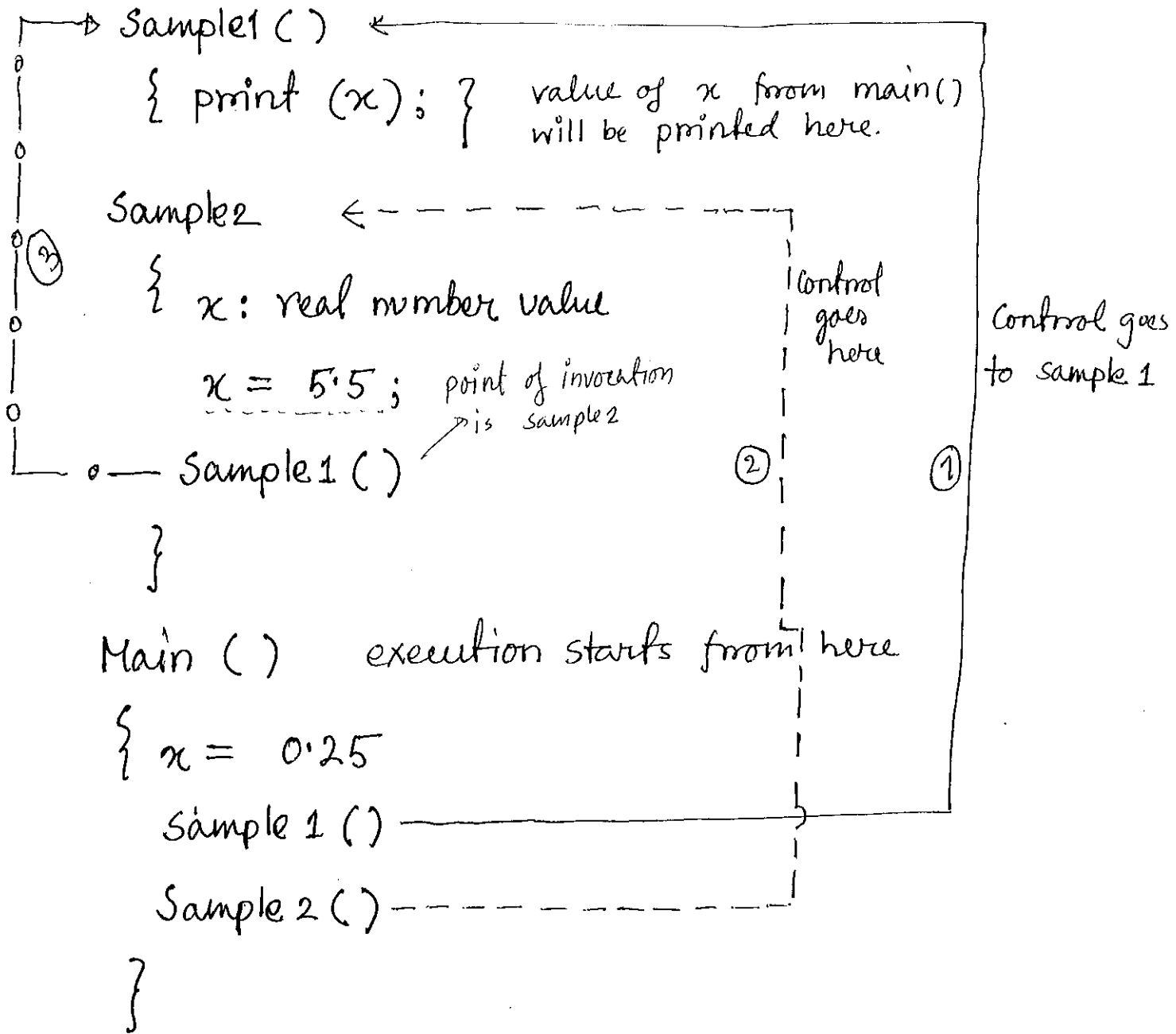
suppose, in main(), we have
main() {x=10;}
it will print x=10;

* Main fn^c is often known as ancestor block

Example question on static and dynamic scoping :

Pseudocode:

x : real number value



Static 0.25
0.25

Dynamic 0.25
5.5

Example question: Hypothetical language.

```
int i  
program main ()  
{ i = 10;  
  call fn ();  
}
```

```
procedure fn ()  
{  
  int i = 20; point of invocation.  
  call g ();  
}
```

```
procedure g ()  
{  
  print (i);  
}
```

Static: $x = 10$ \rightarrow global value

Dynamic: $y = 20$

ALGORITHMIC LANGUAGE ALGOL 60

∗ ALGOL 60 was the outcome of a combined quest for the search of an universal programming language.



In fact, it is the

First step toward sophistication

Motivation:

- ∗ FORTRAN arrived, but it was for IBM
- ∗ Other languages were developed, but they were for specific machines.
- ∗ No portable language.
 - ↳ All languages developed till then were machine specific.
- ∗ Absence of a universal programming language for communicating algorithms.

Goals of the language:

- ∗ Close to mathematical notation
syntax of the language.
- ∗ Should be good for algorithms
- ∗ It should be translatable to machine code.