



NORTH-SOUTH UNIVERSITY

Home-Work

Topic: Grammars & Derivations

Course Code: CSE-425

Sec: 1

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BNF: Backus-Naur Form

- A notation technique for Context-Free Grammar.
- Used to describe the syntax of languages used in computing.
- Have 2 sides, Left-Hand Side & Right-Hand Side
- In Left-Hand Side, there can only be Non-Terminals.
- In Right-Hand Side, there can be Non-Terminals/ Lexemes / Tokens.
- L.H.S + R.H.S = Grammar Rule

Example

<program> → begin <stmt_list> end

<stmt_list> → <stmt> | <stmt>; <stmt_list>

<stmt> → <var> = <expression>

<expression> → <var> + <var> | <var> - <var> | <var>

Derivations

A derivation is basically a sequence of production rules, in order to get the input string. During parsing, we take two decisions for some sentential form of input:

- Deciding the non-terminal which is to be replaced.
- Deciding the production rule, by which, the non-terminal will be replaced.

Example

begin A = B + C; B = C end

<program> → begin <stmt_list> end

begin <stmt>; <stmt_list> end

begin <var> = <expression>; <stmt_list> end

begin A = <expression>; <stmt_list> end

begin A = <var> + <var>; <stmt_list> end

begin A = B + <var>; <stmt_list> end

begin A = B + C; <stmt_list> end

begin A = B + C; <stmt> end

begin A = B + C; <var> = <expression> end

begin A = B + C; B = <expression> end

begin A = B + C; B = <var> end

begin A = B + C; B = C end

To decide which non-terminal to be replaced with production rule, we can have two options:

Left-most Derivation

If the sentential form of an input is scanned and replaced from left to right, it is called left-most derivation. The sentential form derived by the left-most derivation is called the left-sentential form.

Right-most Derivation

If we scan and replace the input with production rules, from right to left, it is known as right-most derivation. The sentential form derived from the right-most derivation is called the right-sentential form.

Sample Grammar

<exp> → <exp> + <exp> | <exp> * <exp> | <exp> | <number>

<number> → <number> <digit> | <digit>

<digit> → 0 | 1 | 2 | 3 | 4

Show 234 using Left-Most Derivation and Right-Most Derivation.

Left-Most Derivation:

<number> → <number> <digit>
→ <number> <digit> <digit>
→ <digit> <digit> <digit>
→ 2 <digit> <digit>
→ 23 <digit>
→ 234

Right-Most Derivation:

$\langle \text{number} \rangle \rightarrow \langle \text{number} \rangle \langle \text{digit} \rangle$
 $\rightarrow \langle \text{number} \rangle 4$
 $\rightarrow \langle \text{number} \rangle \langle \text{digit} \rangle 4$
 $\rightarrow \langle \text{number} \rangle 34$
 $\rightarrow \langle \text{digit} \rangle 34$
 $\rightarrow 234$

Practice Grammar

- $\langle \text{assgin} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\langle \text{id} \rangle \rightarrow A \mid B \mid C$
 $\langle \text{expr} \rangle \rightarrow \langle \text{id} \rangle + \langle \text{expr} \rangle \mid \langle \text{id} \rangle * \langle \text{expr} \rangle \mid (\langle \text{expr} \rangle) \mid \langle \text{id} \rangle$

Show $A = B * (A+C)$ using Left-Most Derivation.

Solution:

$\langle \text{assgin} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\rightarrow A = \langle \text{expr} \rangle$
 $\rightarrow A = \langle \text{id} \rangle * \langle \text{expr} \rangle$
 $\rightarrow A = B * \langle \text{expr} \rangle$
 $\rightarrow A = B * (\langle \text{expr} \rangle)$
 $\rightarrow A = B * (\langle \text{id} \rangle + \langle \text{expr} \rangle)$
 $\rightarrow A = B * (A + \langle \text{expr} \rangle)$
 $\rightarrow A = B * (A + \langle \text{id} \rangle)$
 $\rightarrow A = B * (A + C)$

Show $A = A * (B + (C * A))$ using Left-Most Derivation.

Solution:

$\langle \text{assgin} \rangle \rightarrow \langle \text{id} \rangle = \langle \text{expr} \rangle$
 $\rightarrow A = \langle \text{expr} \rangle$
 $\rightarrow A = \langle \text{id} \rangle * \langle \text{expr} \rangle$
 $\rightarrow A = A * \langle \text{expr} \rangle$
 $\rightarrow A = A * (\langle \text{expr} \rangle)$
 $\rightarrow A = A * (\langle \text{id} \rangle + \langle \text{expr} \rangle)$
 $\rightarrow A = A * (B + \langle \text{expr} \rangle)$
 $\rightarrow A = A * (B + (\langle \text{expr} \rangle))$
 $\rightarrow A = A * (B + (\langle \text{id} \rangle * \langle \text{expr} \rangle))$
 $\rightarrow A = A * (B + (C * \langle \text{expr} \rangle))$
 $\rightarrow A = A * (B + (C * \langle \text{id} \rangle))$
 $\rightarrow A = A * (B + (C * A))$

2. Roll $\rightarrow \langle \text{ROLL} \rangle$ Class Studs $\langle / \text{ROLL} \rangle$
Class $\rightarrow \langle \text{CLASS} \rangle$ Text $\langle / \text{CLASS} \rangle$
Text \rightarrow Char Text
Text \rightarrow Char
Char \rightarrow a... (other chars)
Studs \rightarrow Stud Studs | Stud
Stud $\rightarrow \langle \text{STUD} \rangle$ TEXT $\langle / \text{STUD} \rangle$

Generate documents such as: $\langle \text{ROLL} \rangle \langle \text{CLASS} \rangle \text{cs154} \langle / \text{CLASS} \rangle$

$\langle \text{STUD} \rangle \text{Sally} \langle / \text{STUD} \rangle$

$\langle \text{STUD} \rangle \text{Fred} \langle / \text{STUD} \rangle$

$\langle / \text{ROLL} \rangle$

Try it yourself to practice.

